# Network Science 

## Graphs and Networks

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## The Bridges of Konigsberg

## Drawing Curves with a Single Stroke...



# Königsberg (today's Kaliningrad, Russia) 



## Konigsberg's People



Immanuel Kant (1724-1804)


Gustav Kirchhoff
(1824-1887)


David Hilbert (1862-1943)

## THE BRIDGES OF KONIGSBERG



Can one walk across the seven bridges and never cross the same bridge twice?

## THE BRIDGES OF KONIGSBERG



Can one walk across the seven bridges and never cross the same bridge twice?

1735: Euler's theorem:
(a) If a graph has more than two nodes of odd degree, there is no path.
(b) If a graph is connected and has no odd degree nodes, or two such vertices, it has at least one path.

Euler's solution is considered to be the first theorem in graph theory.

## The Bridges Today



## A "Local" Variation of Euler's Problem



## Graphs and networks after the "bridges"

- Laws of electrical circuitry (G. Kirchhoff, 1845)
- Molecular structure in chemistry (A. Cayley, 1874)
- Network representation of social interactions (J. Moreno, 1930)
- Power grids (1910)
- Telecommunications and the Internet (1960)
- Google (1997), Facebook (2004), Twitter (2006), . . .


## Networks and graphs

## COMPONENTS OF A COMPLEX SYSTEM



## NETWORKS OR GRAPHS?

network often refers to real systems
-www,
-social network
-metabolic network.
Language: (Network, node, link)
graph: mathematical representation of a network
-web graph,
-social graph (a Facebook term)
Language: (Graph, vertex, edge)
We will try to make this distinction whenever it is appropriate, but in most cases we will use the two terms interchangeably.

## A COMMON LANGUAGE



## CHOOSING A PROPER REPRESENTATION

The choice of the proper network representation determines our ability to use network theory successfully.

In some cases there is a unique, unambiguous representation. In other cases, the representation is by no means unique.

For example, the way we assign the links between a group of individuals will determine the nature of the question we can study.

## CHOOSING A PROPER REPRESENTATION



## CHOOSING A PROPER REPRESENTATION



If you connect individuals based on their first name (all Peters connected to each other), you will be exploring what?

It is a network, nevertheless.

## UNDIRECTED VS. DIRECTED NETWORKS

## Undirected

Links: undirected (symmetrical)
Graph:


Undirected links:
coauthorship links
Actor network
protein interactions

## Directed

Links: directed (arcs).
Digraph = directed graph:


An undirected link is the superposition of two opposite directed links.

Directed links:
URLs on the www
phone calls
metabolic reactions

## Section 2.2

## Reference Networks

| NETWORK | NODES | LINKS | DIRECTED UNDIRECTED | N | L |
| :---: | :---: | :---: | :---: | :---: | :---: |
| Internet | Routers | Internet connections | Undirected | 192,244 | 609,066 |
| www | Webpages | Links | Directed | 325,729 | 1,497,134 |
| Power Grid | Power plants, transformers | Cables | Undirected | 4,941 | 6,594 |
| Mobile Phone Calls | Subscribers | Calls | Directed | 36,595 | 91,826 |
| Email | Email addresses | Emails | Directed | 57,194 | 103,731 |
| Science Collaboration | Scientists | Co-authorship | Undirected | 23,133 | 93,439 |
| Actor Network | Actors | Co-acting | Undirected | 702,388 | 29,397,908 |
| Citation Network | Paper | Citations | Directed | 449,673 | 4,689,479 |
| E. Coli Metabolism | Metabolites | Chemical reactions | Directed | 1,039 | 5,802 |
| Protein Interactions | Proteins | Binding interactions | Undirected | 2,018 | 2,930 |

## Degree, Average Degree and Degree Distribution

## NODE DEGREES

Node degree: the number of links connected to the node.
(2)


In directed networks we can define an in-degree and out-degree. The (total) degree is the sum of in- and out-degree.

$$
k_{C}^{\text {in }}=2 \quad k_{C}^{\text {out }}=1 \quad k_{C}=3
$$

Source: a node with $\mathrm{k}^{\mathrm{in}}=0$; Sink: a node with $\mathrm{k}^{\text {out }}=0$.

## A BIT OF STATISTICS

## BRIEF STATISTICS REVIEW

Four key quantities characterize a sample of $N$ values $X_{1}, \ldots, X_{N}$ :

Average (mean):
$\langle x\rangle=\frac{x_{1}+x_{2}+\ldots+x_{N}}{N}=\frac{1}{N} \sum_{i=1}^{N} x_{i}$

The $n^{\text {th }}$ moment:
$\left\langle x^{n}\right\rangle=\frac{x_{1}^{n}+x_{2}^{n}+\ldots+x_{N}^{n}}{N}=\frac{1}{N} \sum_{i=1}^{N} x_{i}^{n}$

## Standard deviation:

$$
\sigma_{x}=\sqrt{\frac{1}{N} \sum_{i=1}^{N}\left(x_{i}-\langle x\rangle\right)^{2}}
$$

Distribution of $x$ :

$$
p_{x}=\frac{1}{N} \sum_{i} \delta_{x, x_{i}}
$$

where $p_{x}$ follows

$$
\sum_{i} p_{x}=1\left(\int p_{x} d x=1\right)
$$

## AVERAGE DEGREE



$$
\langle k\rangle \equiv \frac{1}{N} \sum_{i=1}^{N} k_{i} \quad\langle k\rangle \equiv \frac{2 L}{N}
$$

N - the number of nodes in the graph


## Average Degree

| NETWORK | NODES | LINKS | DIRECTED <br> UNDIRECTED | N | L | $\langle k\rangle$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Internet | Routers | Internet connections | Undirected | 192,244 | 609,066 | 6.33 |
| www | Webpages | Links | Directed | 325,729 | 1,497,134 | 4.60 |
| Power Grid | Power plants, transformers | Cables | Undirected | 4.941 | 6,594 | 2.67 |
| Mobile Phone Calls | Subscribers | Calls | Directed | 36,595 | 91,826 | 2.51 |
| Email | Email addresses | Emails | Directed | 57,194 | 103,731 | 1.81 |
| Science Collaboration | Scientists | Co-authorship | Undirected | 23,133 | 93,439 | 8.08 |
| Actor Network | Actors | Co-acting | Undirected | 702,388 | 29,397,908 | 83.71 |
| Citation Network | Paper | Citations | Directed | 449,673 | 4,689,479 | 10.43 |
| E. Coli Metabolism | Metabolites | Chemical reactions | Directed | 1,039 | 5,802 | 5.58 |
| Protein Interactions | Proteins | Binding interactions | Undirected | 2,018 | 2,930 | 2.90 |

## DEGREE DISTRIBUTION

## Degree distribution

$P(k)$ : probability that a randomly chosen node has degree $k$
$\mathrm{N}_{\mathrm{k}}=$ \# nodes with degree k
$P(k)=N_{k} / N \rightarrow$ plot
c.

b.



## DEGREE DISTRIBUTION




Image 2.4b


## Adjacency matrix


$\mathbf{A}_{\mathrm{ij}}=\mathbf{1}$ if there is a link between node $i$ and $j$
$\mathbf{A}_{\mathrm{ij}}=\mathbf{0}$ if nodes $i$ and $j$ are not connected to each other.

$$
A_{i j}=\left(\begin{array}{cccc}
0 & 1 & 0 & 1 \\
1 & 0 & 0 & 1 \\
0 & 0 & 0 & 1 \\
1 & 1 & 1 & 0
\end{array}\right) \quad A_{i j}=\left(\begin{array}{cccc}
0 & 0 & 0 & 0 \\
1 & 0 & 0 & 1 \\
0 & 0 & 0 & 1 \\
1 & 0 & 0 & 0
\end{array}\right)
$$

Note that for a directed graph (right) the matrix is not symmetric.
$A_{i j}=1$ if there is a link pointing from node $j$ and $i$
$A_{i j}=0$ if there is no link pointing from $j$ to $i$.

## ADJACENCY MATRIX AND NODE DEGREES

$$
\begin{array}{rlr}
A_{i j}=\left(\begin{array}{llll}
0 & 1 & 0 & 1 \\
1 & 0 & 0 & 1 \\
0 & 0 & 0 & 1 \\
1 & 1 & 1 & 0
\end{array}\right) & k_{i}=\sum_{i=1}^{N} A_{i j} \\
A_{i j}=A_{j i} & k_{j}=\sum_{i=1}^{N} A_{i j} \\
A_{i i} & =0 & L=\frac{1}{2} \sum_{i=1}^{N} k_{i}=\frac{1}{2} \sum_{i j}^{N} A_{i j}
\end{array}
$$

Directed


$$
A_{i j}=\left(\begin{array}{ccc|c}
0 & 0 & 0 & 0 \\
1 & 0 & 0 & 1 \\
0 & 0 & 0 & 1 \\
1 & 0 & 0 & 0
\end{array}\right)
$$

$$
k_{i}^{i n}=\sum_{j=1}^{N} A_{i j}
$$

$$
k_{j}^{\text {out }}=\sum_{i=1}^{N} A_{i j}
$$

$$
L=\sum_{i=1}^{N} k_{i}^{\text {in }}=\sum_{j=1}^{N} k_{j}^{\text {out }}=\sum_{i, j}^{N} A_{i j}
$$

## ADJACENCY MATRIX



## Real networks are sparse

## COMPLETE GRAPH

The maximum number of links a network of N nodes can have is: $L_{\max }=\binom{N}{2}=\frac{N(N-1)}{2}$


A graph with degree $L=L_{\text {max }}$ is called a complete graph, and its average degree is $<\mathrm{k}>=\mathrm{N}-\mathbf{1}$

## Most networks observed in real systems are sparse:

$$
\begin{gathered}
L \ll L_{\max } \\
\quad \text { or } \\
<k>\ll N-1 .
\end{gathered}
$$

| WWW (ND Sample): | $\mathrm{N}=325,729 ;$ | $\mathrm{L}=1.410^{6}$ | $\mathrm{~L}_{\max }=10^{12}$ | $<\mathrm{k}>=4.51$ |
| :--- | :--- | :--- | :--- | :--- |
| Protein (S. Cerevisiae): | $\mathrm{N}=1,870 ;$ | $\mathrm{L}=4,470$ | $\mathrm{~L}_{\max }=10^{7}$ | $<\mathrm{k}>=2.39$ |
| Coauthorship (Math): | $\mathrm{N}=70,975 ;$ | $\mathrm{L}=210^{5}$ | $\mathrm{~L}_{\max }=310^{10}$ | $<\mathrm{k}>=3.9$ |
| Movie Actors: | $\mathrm{N}=212,250 ;$ | $\mathrm{L}=610^{6}$ | $\mathrm{~L}_{\max }=1.810^{13}$ | $<\mathrm{k}>=28.78$ |

(Source: Albert, Barabasi, RMP2002)

## ADJACENCY MATRICES ARE SPARSE

## WEIGHTED AND UNWEIGHTED NETWORKS

$$
A_{i j}=w_{i j}
$$

## BIPARTITE NETWORKS

## BIPARTITE GRAPHS

bipartite graph (or bigraph) is a graph whose nodes can be divided into two disjoint sets $U$ and $V$ such that every link connects a node in $U$ to one in $V$; that is, $U$ and $V$ are independent sets.


## GENE NETWORK - DISEASE NETWORK

Goh, Cusick, Valle, Childs, Vidal \& Barabási, PNAS (2007)


The human diseaseome is a bipartite network, whose nodes are diseases ( $U$ ) and genes $(V)$, in which a disease is connected to a gene if mutations in that gene are known to affect the particular disease

## HUMAN DISEASE NETWORK



## PATHOLOGY

## PATHS

A path is a sequence of nodes in which each node is adjacent to the next one
$P_{i 0, i n}$ of length $n$ between nodes $\mathrm{i}_{0}$ and $\mathrm{i}_{n}$ is an ordered collection of $n+1$ nodes and $n$ links

$$
P_{n}=\left\{i_{0}, i_{1}, i_{2}, \ldots, i_{n}\right\} \quad P_{n}=\left\{\left(i_{0}, i_{1}\right),\left(i_{1}, i_{2}\right),\left(i_{2}, i_{3}\right), \ldots,\left(i_{n-1}, i_{n}\right)\right\}
$$



- In a directed network, the path can follow only the direction of an arrow.


The distance (shortest path, geodesic path) between two nodes is defined as the number of edges along the shortest path connecting them.
*If the two nodes are disconnected, the distance is infinity.

In directed graphs each path needs to follow the direction of the arrows.
Thus in a digraph the distance from node $A$ to $B$ (on an $A B$ path) is generally different from the distance from node $B$ to $A$ (on a BCA path).

## $\mathbf{N}_{\mathrm{ij}}$, number of paths between any two nodes $\boldsymbol{i}$ and $\boldsymbol{j}$ :

Length $n=1$ : If there is a link between $i$ and $j$, then $\mathrm{A}_{\mathrm{ij}}=1$ and $\mathrm{A}_{\mathrm{ij}}=0$ otherwise.
Length n=2: If there is a path of length two between $i$ and $j$, then $\mathrm{A}_{\mathrm{ik}} \mathrm{A}_{\mathrm{kj}}=1$, and $\mathrm{A}_{\mathrm{ik}} \mathrm{A}_{\mathrm{kj}}=0$ otherwise. The number of paths of length 2 :

$$
N_{i j}^{(2)}=\sum_{k=1}^{N} A_{i k} A_{k j}=\left[A^{2}\right]_{i j}
$$

Length n: In general, the number of paths of length $n$ between $i$ and $j$ is*

$$
N_{i j}^{(n)}=\left[A^{n}\right]_{i j}
$$

[^0]
## FINDING DISTANCES: BREADTH FIRST SEARCH

Distance between node 0 and node 4:
1.Start at 0 .


## FINDING DISTANCES: BREADTH FIRST SEARCH

Distance between node 0 and node 4:

## 1. Start at 0 .

2. Find the nodes adjacent to 1 . Mark them as at distance 1. Put them in a queue.


## FINDING DISTANCES: BREADTH FIRST SEARCH

## Distance between node 0 and node 4:

1. Start at 0.
2. Find the nodes adjacent to 0 . Mark them as at distance 1 . Put them in a queue.
3. Take the first node out of the queue. Find the unmarked nodes adjacent to it in the graph. Mark them with the label of 2 . Put them in the queue.


## FINDING DISTANCES: BREADTH FIRST SEARCH

## Distance between node 0 and node 4:

1.Repeat until you find node 4 or there are no more nodes in the queue.
2. The distance between 0 and 4 is the label of 4 or, if 4 does not have a label, infinity.


## FINDING DISTANCES: BREADTH FIRST SEARCH



The computational complexity of the BFS algorithm, representing the approximate number of steps the computer needs to find dij on a network of N nodes and L links, is $\mathrm{O}(\mathrm{N}+\mathrm{L})$

BFS can be used to test bipartiteness, by starting the search at any vertex and giving alternating labels to the vertices visited during the search.

That is, give label 0 to the starting vertex, 1 to all its neighbors, 0 to those neighbors' neighbors, and so on.

If at any step a vertex has (visited) neighbors with the same label as itself, then the graph is not bipartite.

If the search ends without such a situation occurring, then the graph is bipartite.

Try here!

A) A Bipartite Graph

B) A non-Bipartite Graph

## Note

A graph is bipartite iff it contains no odd cycle.

## NETWORK DIAMETER AND AVERAGE DISTANCE

Diameter ( $d_{\text {max }}$ ): the maximum distance between any pair of nodes in the graph.

$$
d_{\max } \equiv \max _{i \neq j} d_{i j}
$$

where $d_{i j}$ is the distance from node $i$ to node $j$

Average distance ( <d> ): for a connected graph

$$
\langle d\rangle \equiv \frac{1}{N(N-1)} \sum_{i} \sum_{j \neq i} d_{i j}
$$

## Shortest Path



The path with the shortest length between two nodes (distance).

## Diameter

## Average Path Length



The longest shortest path in a graph


The average of the shortest paths for all pairs of nodes.


A path with the same start and end node.


A path that traverses each link exactly once.

## Hamiltonian Path



A path that visits each node exactly once.

## CONNECTEDNESS

## CONNECTIVITY OF UNDIRECTED GRAPHS

Connected (undirected) graph: any two vertices can be joined by a path. A disconnected graph is made up by two or more connected components.


Largest Component: Giant Component

The rest: Isolates

Bridge: if we erase it, the graph becomes disconnected.

## CONNECTIVITY OF UNDIRECTED GRAPHS

The adjacency matrix of a network with several components can be written in a blockdiagonal form, so that nonzero elements are confined to squares, with all other elements being zero:
a.

b.


## CONNECTIVITY OF DIRECTED GRAPHS

Strongly connected directed graph: has a path from each node to every other node and vice versa (e.g. AB path and BA path).

Weakly connected directed graph: it is connected if we disregard the edge directions.


## Section 2.9

Box 2.6

## Finding the Connected Components of a Network

- Start from a randomly chosen node $i$ and perform a BFS (BOX 2.5). Label all nodes reached this way with $n=1$.
- If the total number of labeled nodes equals $N$, then the network is connected. If the number of labeled nodes is smaller than $N$, the network consists of several components. To identify them, proceed to step 3.
- Increase the label $n \rightarrow n+1$. Choose an unmarked node $j$, label it with $n$. Use BFS to find all nodes reachable from $j$, label them all with $n$. Return to step 2 .


## Clustering coefficient and cliques

## CLUSTERING COEFFICIENT

What fraction of your neighbors are connected?

* Node i with degree $k_{i}$
* $e_{i}=$ number of links between the $k_{i}$ neighbors of $i$

$$
C_{i}=\frac{2 e_{i}}{k_{i}\left(k_{i}-1\right)}
$$

Note: $0 \leq C_{i} \leq 1$
a.

$C_{i}=1$

C. $=1 / 2$

$C_{i}=0$
b.

$\langle C\rangle=\frac{13}{42} \approx 0.310$
$C_{\Delta}=\frac{3}{8}=0.375$

## Cliques

Given an unweighted undirected graph $G=(\mathrm{V}, \mathrm{E})$ :

- A clique is a subset of mutually adjacent vertices
- A maximal clique is a clique that is not contained in a larger one
- A maximum clique is a clique having largest cardinality

The clique number, denote $\omega(\mathrm{G})$, is the cardinality of a maximum clique.

Independent set: clique on the complement of G


## summary

## Degree distribution:

$P(k)$

## Path length:

Clustering coefficient:

$$
C_{i}=\frac{2 e_{i}}{k_{i}\left(k_{i}-1\right)}
$$

## GRAPHOLOGY 1

## Undirected



$$
\begin{gathered}
A_{i j}=\left(\begin{array}{llll}
0 & 1 & 1 & 0 \\
1 & 0 & 1 & 1 \\
1 & 1 & 0 & 0 \\
0 & 1 & 0 & 0
\end{array}\right) \\
A_{i i}=0
\end{gathered} \quad A_{i j}=A_{j i} .
$$

Actor network, protein-protein interactions

## Directed

$$
A_{i j}=\left(\begin{array}{llll}
0 & 1 & 0 & 0 \\
0 & 0 & 1 & 1 \\
1 & 0 & 0 & 0 \\
0 & 0 & 0 & 0
\end{array}\right)
$$

$$
A_{i i}=0 \quad A_{i j} \neq A_{j i}
$$

$$
L=\sum_{i, j=1}^{N} A_{i j} \quad\langle k\rangle=\frac{L}{N}
$$

## GRAPHOLOGY <br> 2

## Unweighted (undirected)



$$
A_{i j}=\left(\begin{array}{llll}
0 & 1 & 1 & 0 \\
1 & 0 & 1 & 1 \\
1 & 1 & 0 & 0 \\
0 & 1 & 0 & 0
\end{array}\right)
$$

$$
A_{i i}=0 \quad A_{i j}=A_{j i}
$$

$$
L=\frac{1}{2} \sum_{i, j=1}^{N} A_{i j} \quad<k>=\frac{2 L}{N}
$$

Weighted (undirected)

2

$$
\begin{gathered}
A_{i j}=\left(\begin{array}{cccc}
0 & 2 & 0.5 & 0 \\
2 & 0 & 1 & 4 \\
0.5 & 1 & 0 & 0 \\
0 & 4 & 0 & 0
\end{array}\right) \\
L=\frac{A_{i j}}{2} \sum_{i, j=1}^{N} \text { nonzero }\left(A_{i j}\right) \quad\langle k\rangle=\frac{2 L}{N}
\end{gathered}
$$

Call Graph, metabolic networks

## GRAPHOLOGY

Self-interactions

$$
\begin{aligned}
& A_{i j}=\left(\begin{array}{llll}
1 & 1 & 1 & 0 \\
1 & 0 & 1 & 1 \\
1 & 1 & 0 & 0 \\
0 & 1 & 0 & 1
\end{array}\right) \\
& L=\frac{1}{2} \sum_{i, j=l, i \neq j}^{N} A_{i j}+\sum_{i=1}^{N} A_{i i}
\end{aligned}
$$

Protein interaction network, www

## Multigraph

 (undirected)$$
\begin{aligned}
& A_{i j}=\left(\begin{array}{llll}
0 & 2 & 1 & 0 \\
2 & 0 & 1 & 3 \\
1 & 1 & 0 & 0 \\
0 & 3 & 0 & 0
\end{array}\right) \\
& A_{i i}=0 \\
& L=\frac{1}{2} \sum_{i, j=1}^{N} A_{i j} \quad\langle k\rangle=\frac{2 L}{N}
\end{aligned}
$$

Social networks, collaboration networks

## GRAPHOLOGY 4

## Complete Graph (Clique)

 (undirected)

$$
A_{i j}=\left(\begin{array}{llll}
0 & 1 & 1 & 1 \\
1 & 0 & 1 & 1 \\
1 & 1 & 0 & 1 \\
1 & 1 & 1 & 0
\end{array}\right)
$$

$$
A_{i i}=0 \quad A_{i \neq j}=1
$$

$$
L=L_{\max }=\frac{N(N-1)}{2} \quad<k>=N-1
$$



Actor network, protein-protein interactions

## GRAPHOLOGY: Real networks can have multiple characteristics

## WWW > directed multigraph with self-interactions

Protein Interactions > undirected unweighted with self-interactions

Collaboration network >
undirected multigraph or weighted.

## Mobile phone calls >

directed, weighted.

Facebook Friendship links >
undirected, unweighted.


[^0]:    *holds for both directed and undirected networks.

