



Succinct Games

Describing a game in normal form entails listing all payoffs for all players and strategy combinations. In a game with n players, each facing m pure strategies, one need to store nm^n numbers!

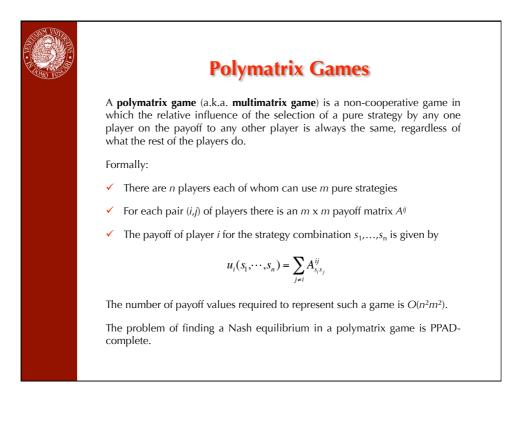
A succinct game (or a succinctly representable game) is a game which may be represented in a size much smaller than its normal-form representation.

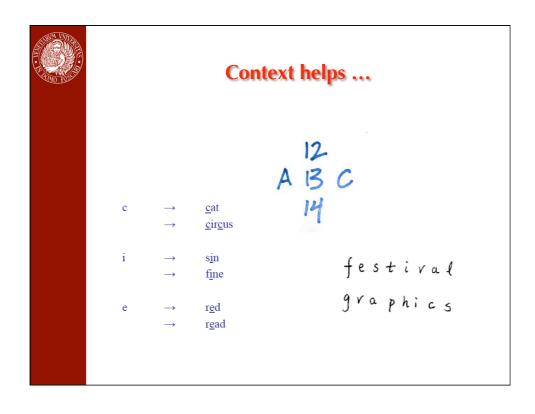
Examples.

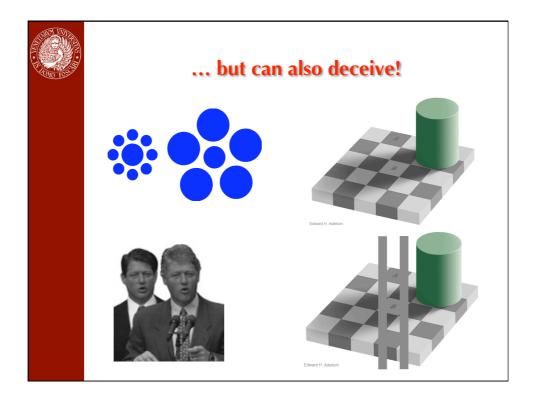
Sparse games. Most of the payoffs are zero.

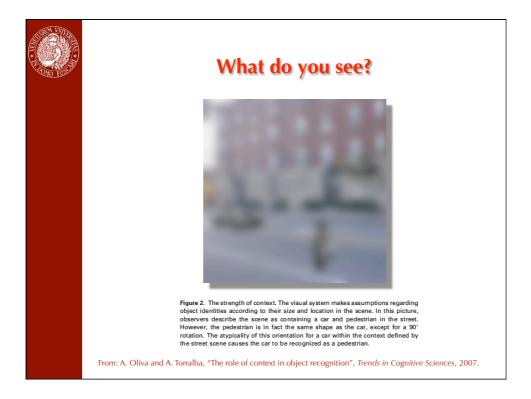
Graphical games. The payoffs of each player depends on the actions of very few (at most *d*) other players. The number of payoffs needed to describe this game is nm^{d+1} .

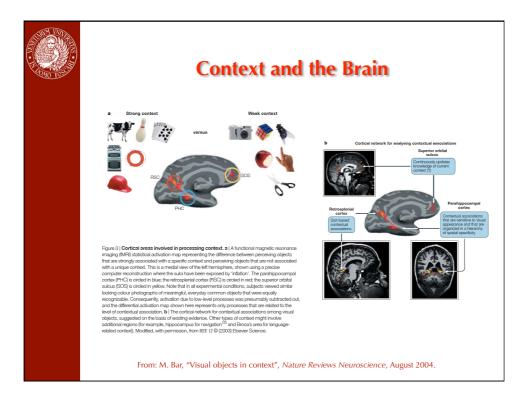
Symmetric games. All players are identical, so in evaluating the payoff of a combination of strategies, all that matters is how many of the *n* players play each of the *s* strategies.

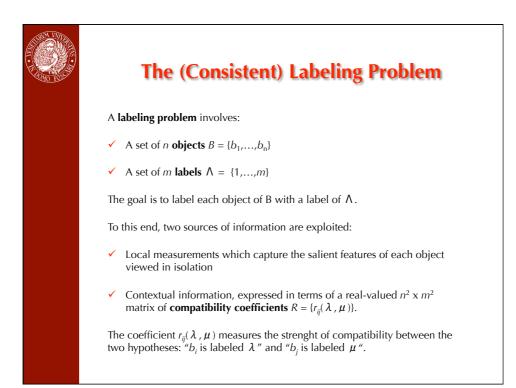












Relaxation Labeling Processes

The initial local measurements are assumed to provide, for each object $b_i \in B$, an *m*-dimensional (probability) vector:

$$p_i^{(0)} = \left(p_i^{(0)}(1), \cdots, p_i^{(0)}(m)\right)^T$$

with $p_i^{(0)}(\lambda) \ge 0$ and $\sum_{\lambda} p_i^{(0)}(\lambda) = 1$. Each $p_i^{(0)}(\lambda)$ represents the initial, non-contextual degree of confidence in the hypothesis " b_i is labeled λ ".

By concatenating vectors $p_1^{(0)}, \dots, p_n^{(0)}$ one obtains an (initial) weighted labeling assignment $p^{(0)} \in \Re^{nm}$.

The space of weighted labeling assignments is

$$IK = \underbrace{\Delta \times \ldots \times \Delta}_{m \text{ times}}$$

where each Δ is the standard simplex of \Re^n . Vertices of IK represent unambiguous labeling assignments

A **relaxation labeling process** takes the initial labeling assignment $p^{(0)}$ as input and iteratively updates it taking into account the compatibility model *R*.



Relaxation Labeling Processes

In a now classic 1976 paper, Rosenfeld, Hummel, and Zucker introduced heuristically the following update rule (assuming a non-negative compatibility matrix):

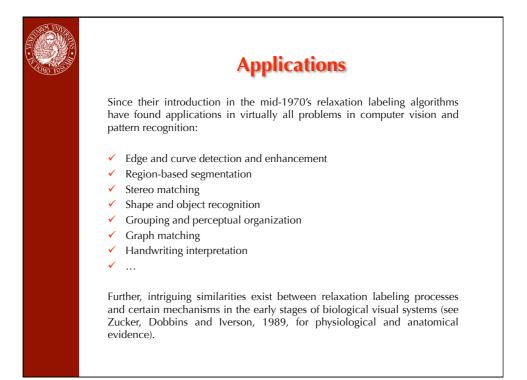
$$p_i^{(t+1)}(\lambda) = \frac{p_i^{(t)}(\lambda)q_i^{(t)}(\lambda)}{\sum_{\mu} p_i^{(t)}(\mu)q_i^{(t)}(\mu)}$$

where

$$q_i^{(t)}(\lambda) = \sum_j \sum_{\mu} r_{ij}(\lambda, \mu) p_i^{(t)}(\mu)$$

quantifies the support that context gives at time *t* to the hypothesis " b_i is labeled with label λ ".

See (Pelillo, 1997) for a rigorous derivation of this rule in the context of a formal theory of consistency.





Hummel and Zucker's Consistency

In 1983, Bob Hummel and Steve Zucker developed an elegant theory of consistency in labeling problem.

By analogy with the unambiguous case, which is easily understood, they define a weighted labeling assignment $p \in IK$ consistent if:

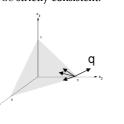
$$\sum_{\lambda} p_i(\lambda) q_i(\lambda) \geq \sum_{\lambda} v_i(\lambda) q_i(\lambda) \qquad i=1 \dots n$$

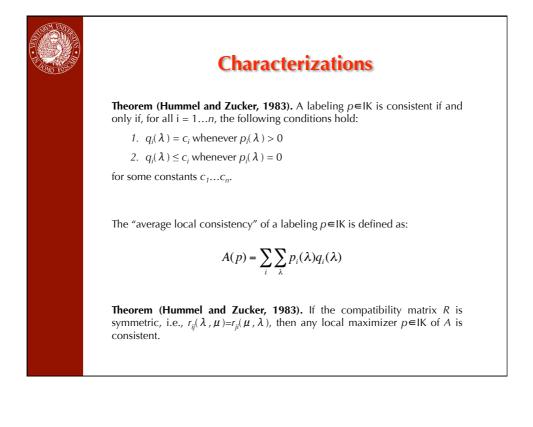
for all labeling assignments $v \in IK$. If strict inequalities hold for all $v \neq p$, then p is said to be **strictly consistent**.

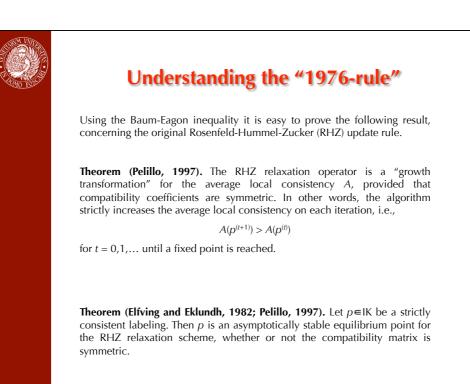
Geometrical interpretation.

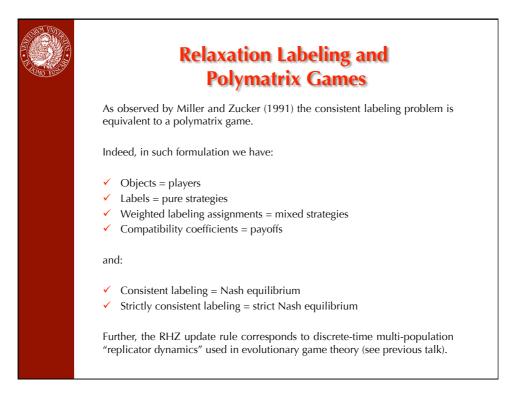
The support vector q points away from all tangent vectors at p (it has null projection in IK).

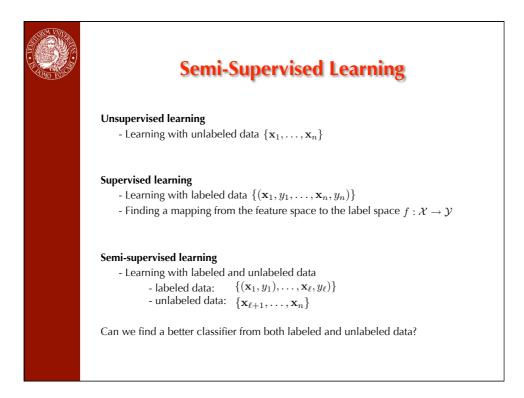
Generalization of classical constraint satisfaction problems!

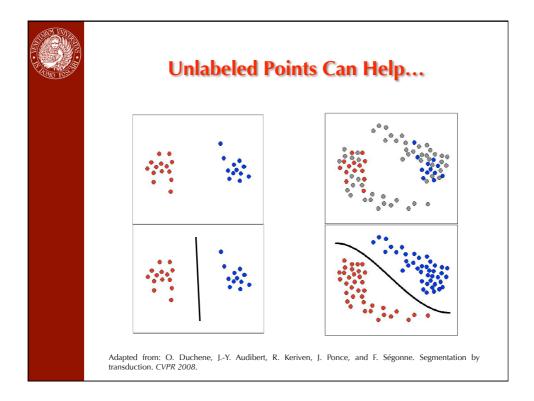


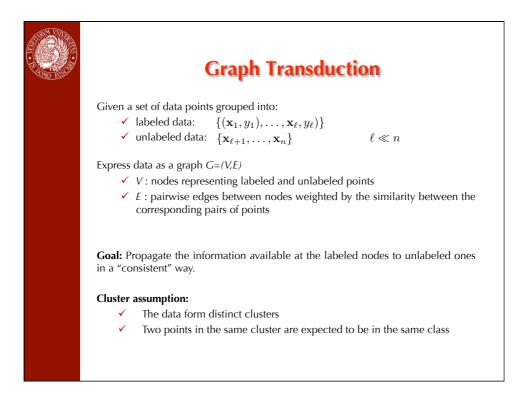


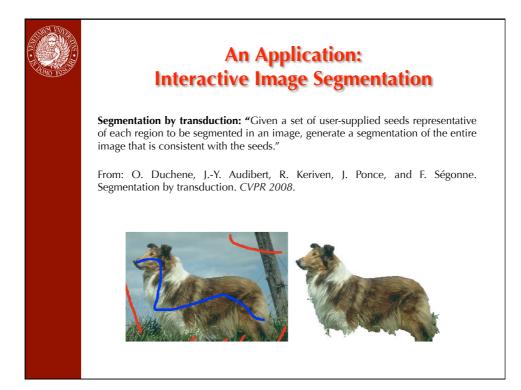


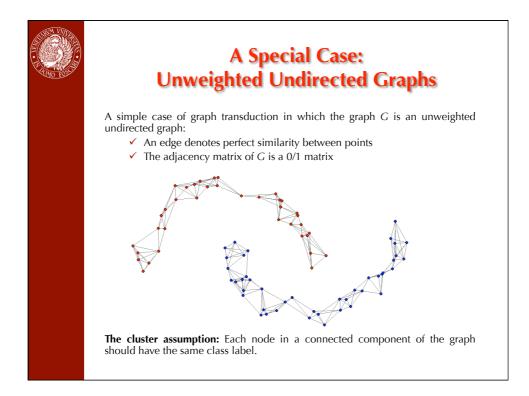


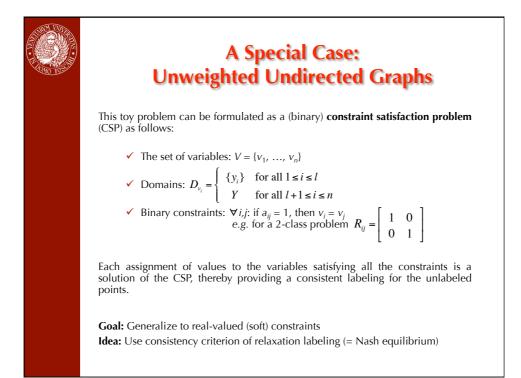


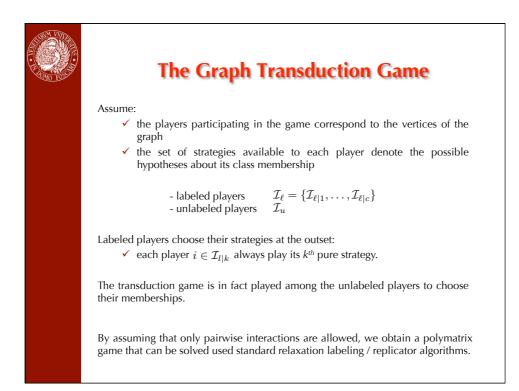


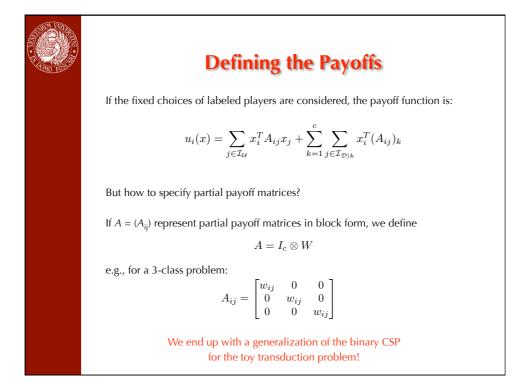












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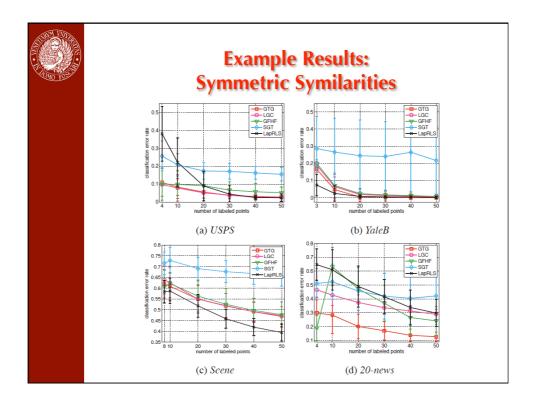
Example Results: Symmetric Symilarities

Data set used: USPS, YaleB, Scene, 20-news

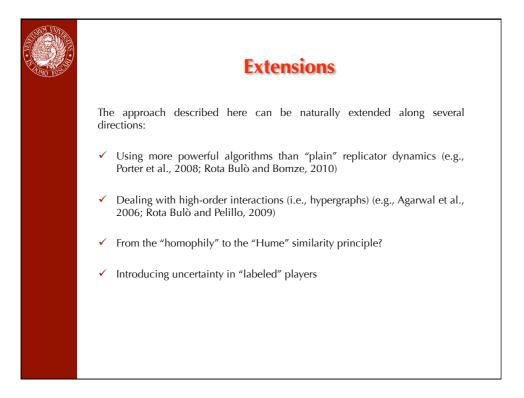
	USPS	YaleB	Scene	20-news
# objects	3874	1755	2688	3970
# dimensions	256	1200	512	8014
# classes	4	3	8	4

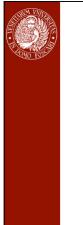
Methods compared:

- ✓ Gaussian fields and harmonic functions (GFHF) (Zhu et al., 2003)
- ✓ Spectral Graph Transducer (SGT) (Joachims, 2003)
- ✓ Local and global consistency (LGC) (Zhou et al., 2004)
- ✓ Laplacian Regularized Least Squares (LapRLS) (Belkin et al., 2006)









References

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